

Digital Data Communication

by

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Digital Transmission



- **Serial Transmission**: data are transferred over a single signal path, and signaling elements are sent down the line one at a time
 - E.g., Device to Device transmission via Wireless/Wired
- **Parallel Transmission**: data are transferred over a parallel set of lines
 - E.g, I/O devices, internal computer signal paths
- Each signaling element may be
 - Less than one bit: with Manchester coding
 - One bit: NRZ-L and FSK are digital and analog examples, respectively
 - More than one bit: QPSK

Transmitter-Receiver Time Synchronization

- Suppose that the sender simply transmits a stream of data bits.
- The sender **has a clock** that governs the timing of the transmitted bits.
- Let, data rate 1Mbps. So, one bit will be transmitted in **every 1 μ s** as measured by the sender's clock.
- the receiver will attempt to sample the medium at the center of each bit time.
- So, sampling would occur once **every 1 μ s**
- Hence, there **will be a problem** if the transmitter's and receiver's clocks are not precisely aligned.
- Let, there is a **drift of 1%**, i.e., the receiver's clock is 1% faster or slower than the transmitter's clock
- Then, the first sampling will be 0.01 of a bit time (i.e., 0.01 μ s) away from the center of the bit (center of bit is 0.5 μ s from beginning and end of bit)
- After 50 or more samples, the receiver may be in error.

Time Synchronization



- Two approaches are common for achieving synchronization
 - asynchronous transmission
 - strategy with this scheme is to avoid the timing problem by not sending long, uninterrupted streams of bits
 - data are transmitted one character at a time, where each character is five to eight bits in length
 - Timing or synchronization must only be maintained within each character
 - synchronous transmission
 - a block of bits is transmitted in a steady stream without start and stop codes
 - The block may be many bits in length
 - To prevent timing drift between transmitter and receiver, their clocks must somehow be synchronized
 - provide a separate clock line for sending clock from sender to receiver
 - embed the clocking information in the data signal. (E.g., Manchester or differential Manchester encoding)

Asynchronous Transmission

The definition of *idle* is equivalent to the signaling element for binary 1.

NRZ-L signaling is common for asynchronous transmission

The **beginning** of a character is signaled by a *start bit* with a value of binary 0

The **final element** is a stop element, which is a binary 1

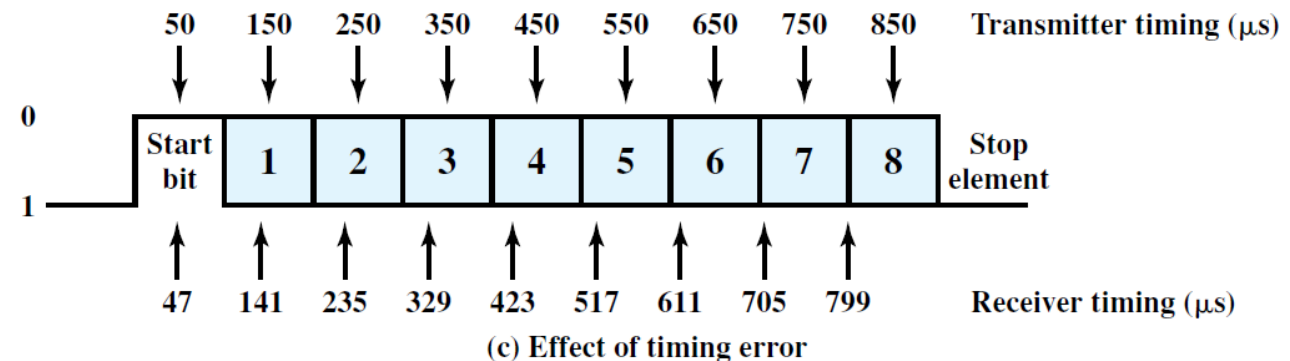
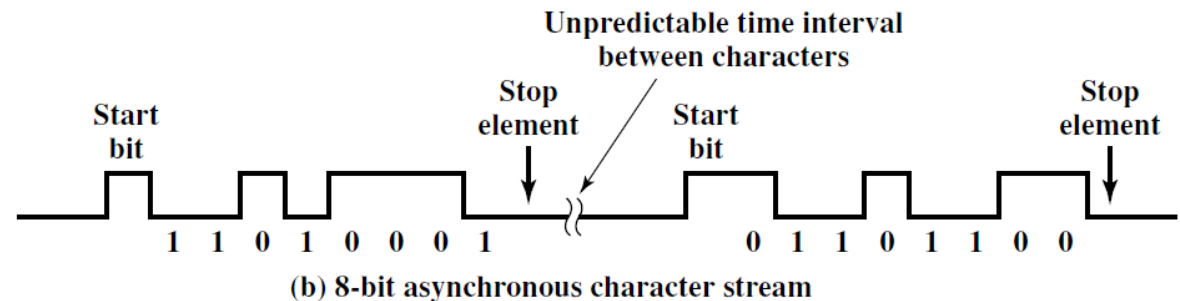
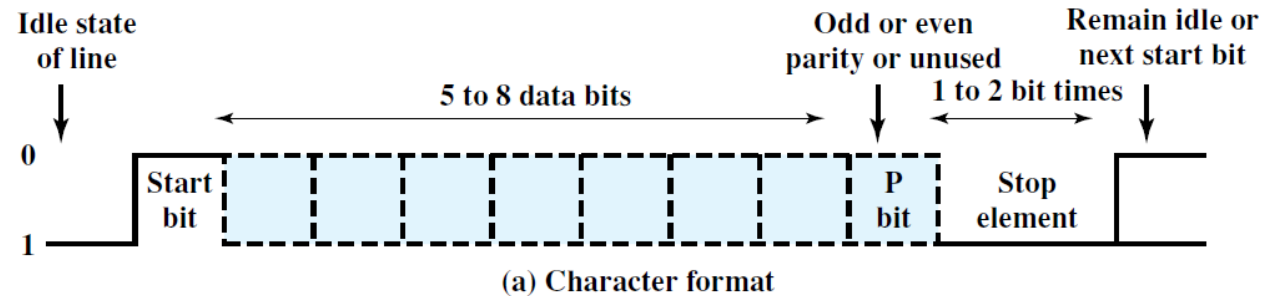


Figure 6.1 Asynchronous Transmission

Cont...

- Asynchronous transmission creates **frame error**.
- **Frame error**: the bit count may be out of alignment due to the cumulative timing error. If bit 7 is a 1 and bit 8 is a 0, bit 8 could be mistaken for a start bit.
- It is simple and cheap but requires an overhead of two to three bits per character

Synchronous Transmission

- Either provide a separate clock line
 - Applicable for short distance communication
- Or, embed the clock information in the data signal
 - Applicable for all
 - E.g. for **Digital signal**: Manchester / Differential Manchester encoding
 - E.g. for **Analog signal**: the carrier frequency itself can be used to synchronize the receiver based on the phase of the carrier.
- With synchronous transmission, there is another level of synchronization required, to allow the receiver to determine the beginning and end of a block of data.
- To achieve this, each block begins with a **preamble** bit pattern and generally ends with a **postamble** bit pattern.
- Asynchronous transmission requires 20% or more overhead. The control information, preamble, and postamble in synchronous transmission are typically less than 100 bits.

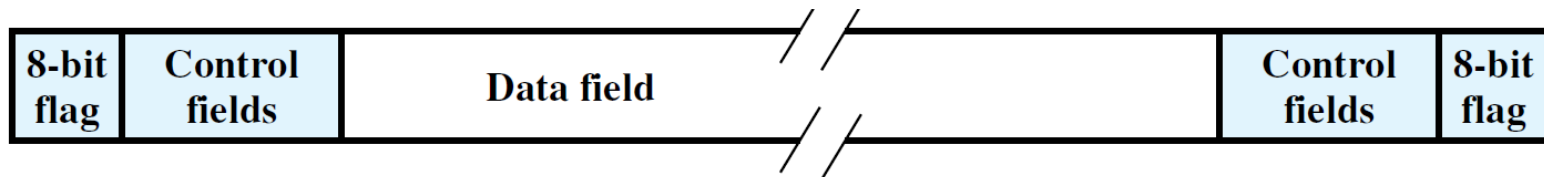


Figure 6.2 Synchronous Frame Format

Line Configurations

- Two characteristics that distinguish various data link configurations
 - Topology
 - half duplex / full duplex transmission
- The **topology** of a data link refers to the physical arrangement of stations on a transmission medium.
- With **half-duplex transmission**, only one of two stations on a point-to-point link may transmit at a time.
- For **full-duplex transmission**, two stations can simultaneously send and receive data from each other.
- With digital signaling, which requires guided transmission, full-duplex operation usually requires two separate transmission paths
- For analog signaling, it depends on frequency

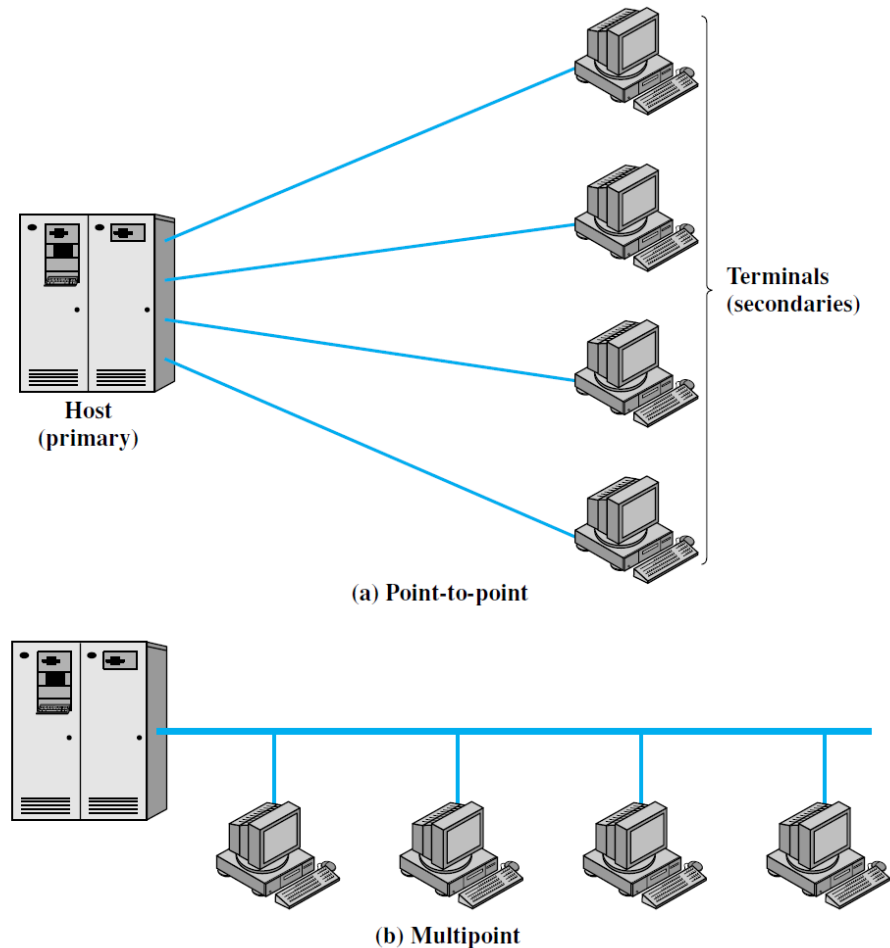
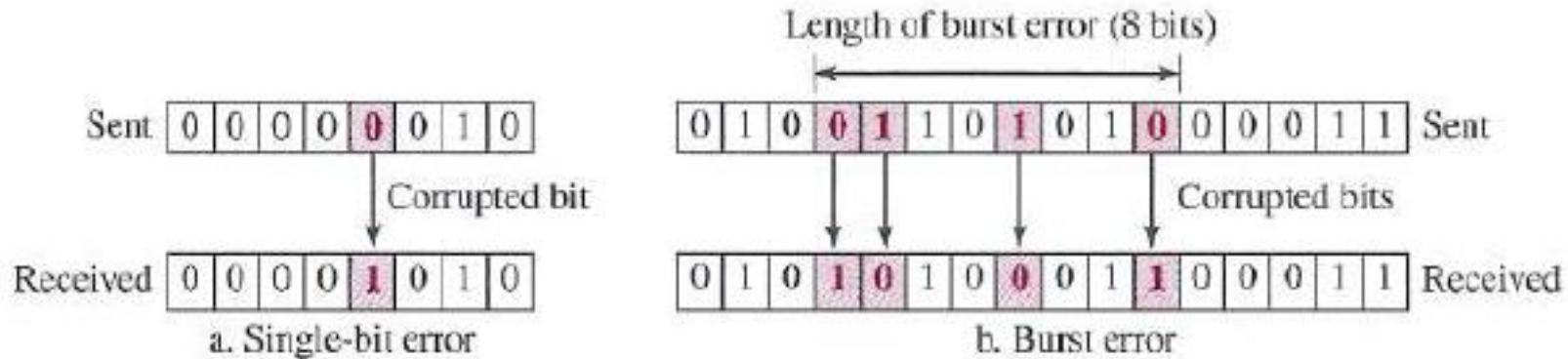


Figure 6.9 Traditional Computer/Terminal Configurations

Types of Error



- Single bit error
- Burst error / multibit error
- **Reason:** noise in the channel

Error Detection and Correction



- Objective:
 - System must guarantee that the data received are identical to the data transmitted
- Two Methods:
 1. If a frame is corrupted between the two nodes, it needs to be corrected (i.e. Detect and Correct error in PHY/MAC layers)
 2. Drop the frame and let the upper layer (e.g. Transport) protocol to handle it by retransmission

Thanks!

Figure and slide materials are taken from the following sources:

1. W. Stallings, (2010), [Data and Computer Communications](#)
2. [NPTL lecture](#) on Data Communication, by Prof. A. K. Pal, IIT Kharagpur
3. B. A. Forouzan, (2013), [Data Communication and Networking](#)