

Internet of Things (IoT)



IEEE 802.15.4

Low-Rate Wireless Networks

: MAC Layer

2011 version: <https://ieeexplore.ieee.org/stamp/stamp.jsp?tp=&arnumber=6012487>

2015 version: <https://ieeexplore.ieee.org/stamp/stamp.jsp?tp=&arnumber=7460875>

2020 version: <https://ieeexplore.ieee.org/stamp/stamp.jsp?tp=&arnumber=9144691>

Dr. Manas Khatua

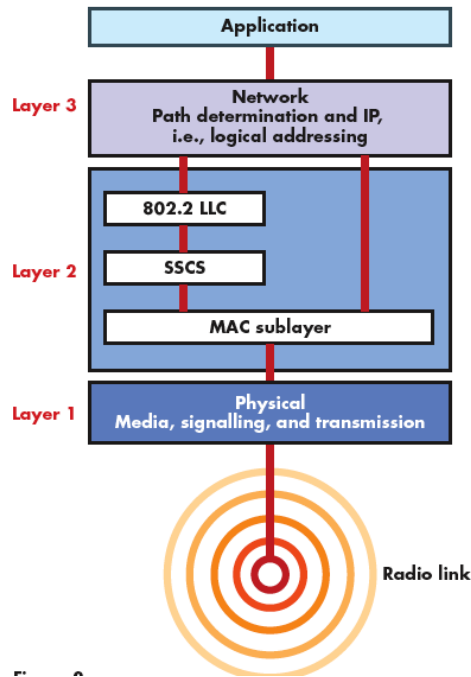
Associate Professor

Dept. of CSE, IIT Guwahati

E-mail: manaskhatua@iitg.ac.in

IEEE 802.15.4 Stack – PHY & MAC

The OSI model adapted to the IEEE 802.15.4



LLC: Logical Link Control – provides protocol multiplexing

SSCS: Service Specific Convergence Sublayer

- IEEE 802.15.4 standard is limited to the **PHY & MAC** Layers
- IEEE 802.15.4 standard MAC provides the **MAC data service** and **MAC management services**.
 - The **MAC data service** enables **transmission** of MAC protocol data units (MPDU) across the PHY data service.
 - The **MAC sublayer features** include
 - beacon management,
 - channel access,
 - GTS management,
 - frame validation,
 - ACK frame delivery,
 - association and disassociation,
 - Device security

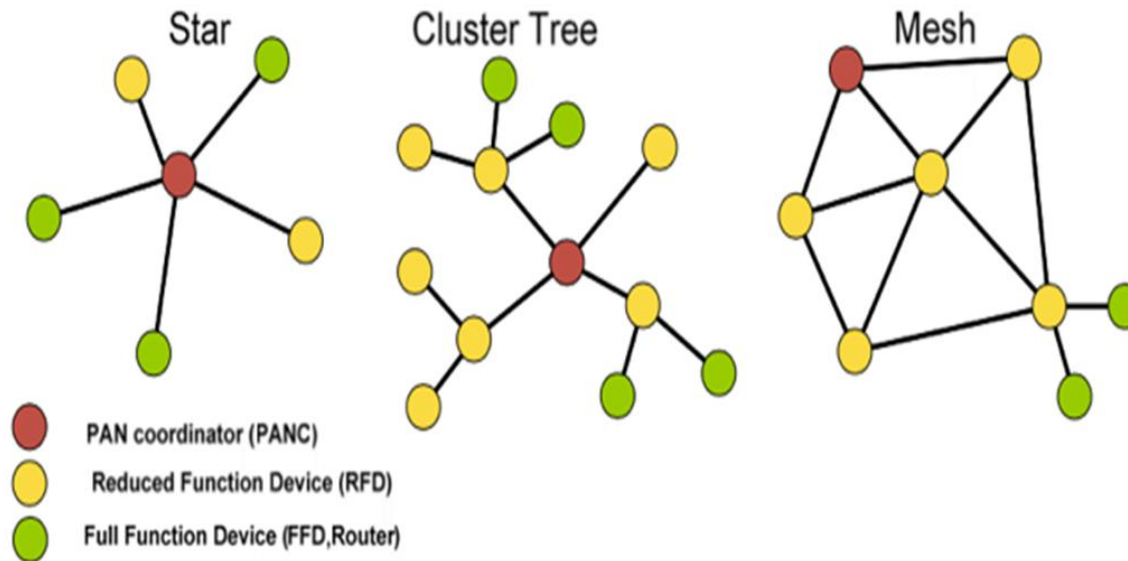
Standard Document: IEEE Std 802.15.4™-2020, “IEEE Standard for **Low-Rate Wireless Networks**”,

Developed by the LAN/MAN Standards Committee of the IEEE Computer Society, Approved on 6 May 2020.

Image Source: <https://www.embedded.com/ieee-802-15-4-zigbee-hardware-and-software-open-the-applications-window/>

IEEE 802.15.4 MAC layer

- MAC layer manages access to the PHY channel
 - defines how devices in the same area will share the frequencies allocated.
- MAC layer establish logical topology of the network

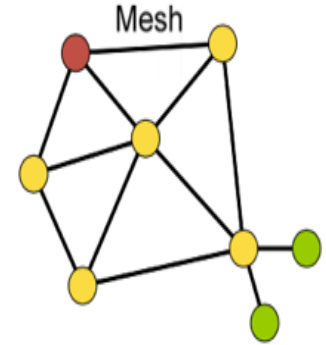


IEEE 802.15.4 Device Types

- There are **two different device types** :
 - full function device (**FFD**)
 - reduced function device (**RFD**)
- The **FFD** can operate in **three modes** by serving as
 - **PAN Coordinator**
 - scanning the network and selecting optimal RF channel
 - selecting the 16 bit PAN ID for the network
 - **Coordinator (aka Parent, Join Proxy)**
 - relaying messages to other FFDs including PAN coordinator
 - transmits periodic beacon (under beacon enable access mode)
 - respond to beacon requests
 - **Device**
 - cannot route messages
 - usually receivers are switched off except during transmission
 - attached to the network only as leaf nodes
- The **RFD** can only serve as:
 - **Device**

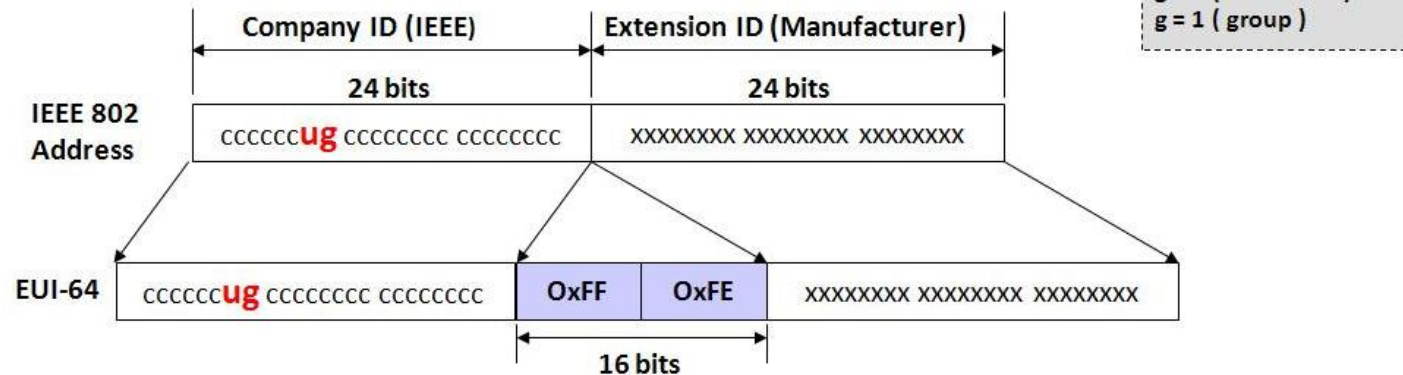
Device Addressing

- Two or more devices **wirelessly** communicating on the **same physical channels** following IEEE 802.15.4 constitute a PAN.
 - A PAN includes at least one FFD (PAN coordinator)
 - Each independent PAN will select a **unique PAN ID**
 - PAN ID is generally of **2 bytes**
- Each device operating on a network has a **unique 64-bit address**
 - called **extended unique identifier** (EUI-64)
 - This address can be used for **direct communication** in the PAN
 - This is also called **64-bit long address**
 - Generally, **it is autogenerated** from unique MAC address
- A device also has a **16-bit short address**, which is **allocated by the PAN coordinator** when the device associates with its coordinator.
 - Note: Same short address may be present into different PAN

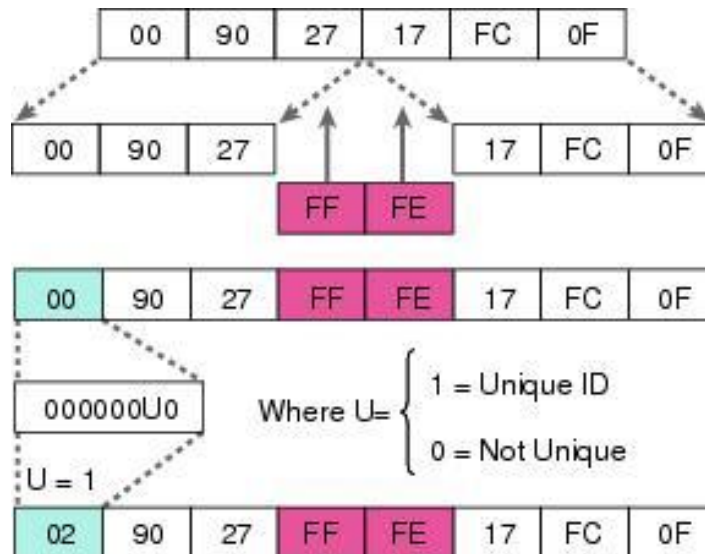


Deriving EUI-64 ID from MAC

Deriving the Modified EUI-64 Interface Identifier from the MAC Address



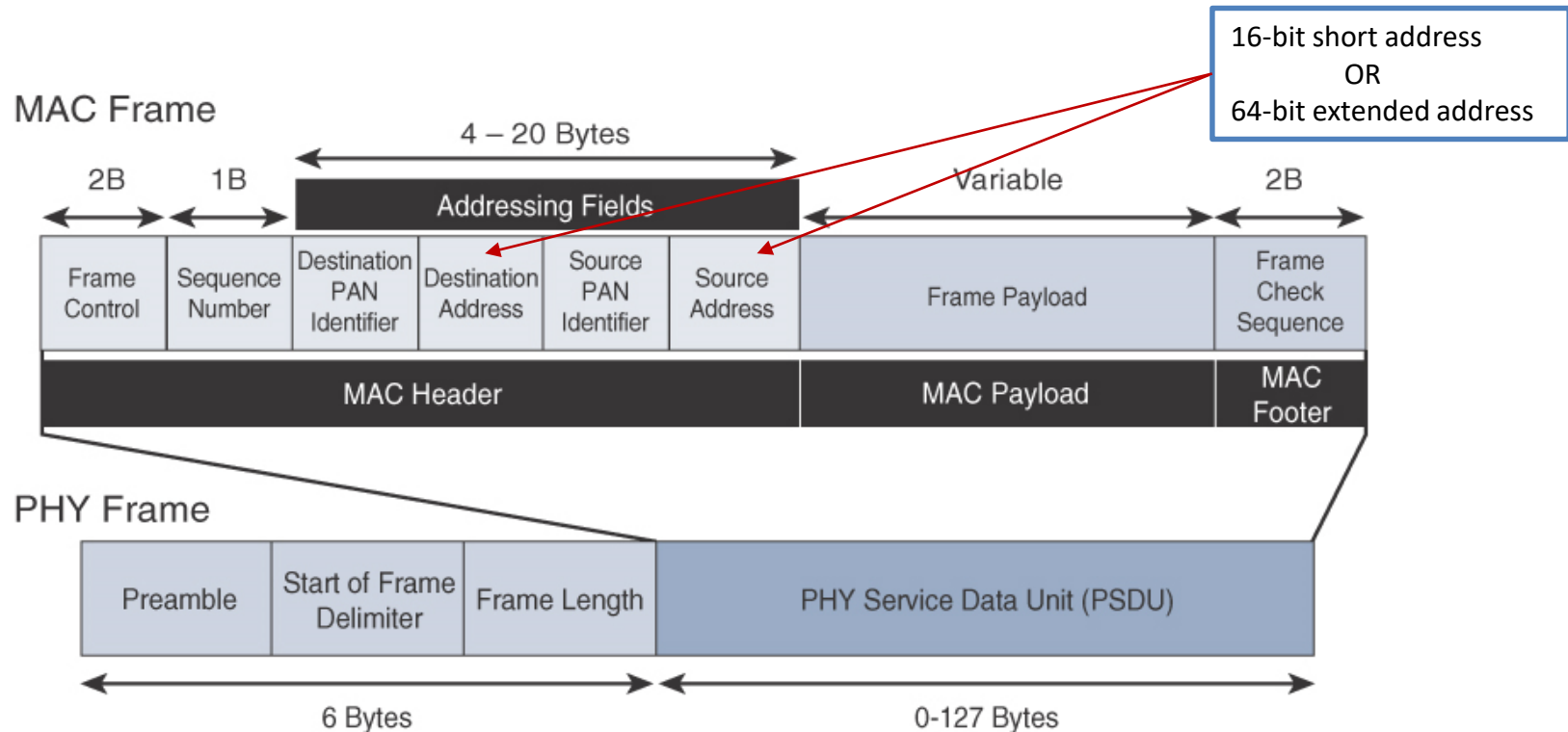
Example:



Addressing Modes

- IEEE 802.15.4 frames contain **address of both the source & destination**.
- **Three different addressing modes**, which sets the address field (none/ short/ long, with/without PAN ID)
 - **Short addressing mode**: The address field includes a short address (2B) & a PAN ID (2B) = (total of 4 bytes).
 - **Long addressing mode**: The address field includes a long address (8B) and a PAN ID (2B) = (total of 10 bytes).
 - **No addressing mode**:
 - For ACK frame - both addresses are missing.
 - For *Data* and *Command* frames - **only one** (either source or destination) field **can be omitted**
 - if the **source address is omitted**, it means the PAN coordinator sent the frame;
 - if the **destination address is missing**, it means it should be received by the PAN coordinator.

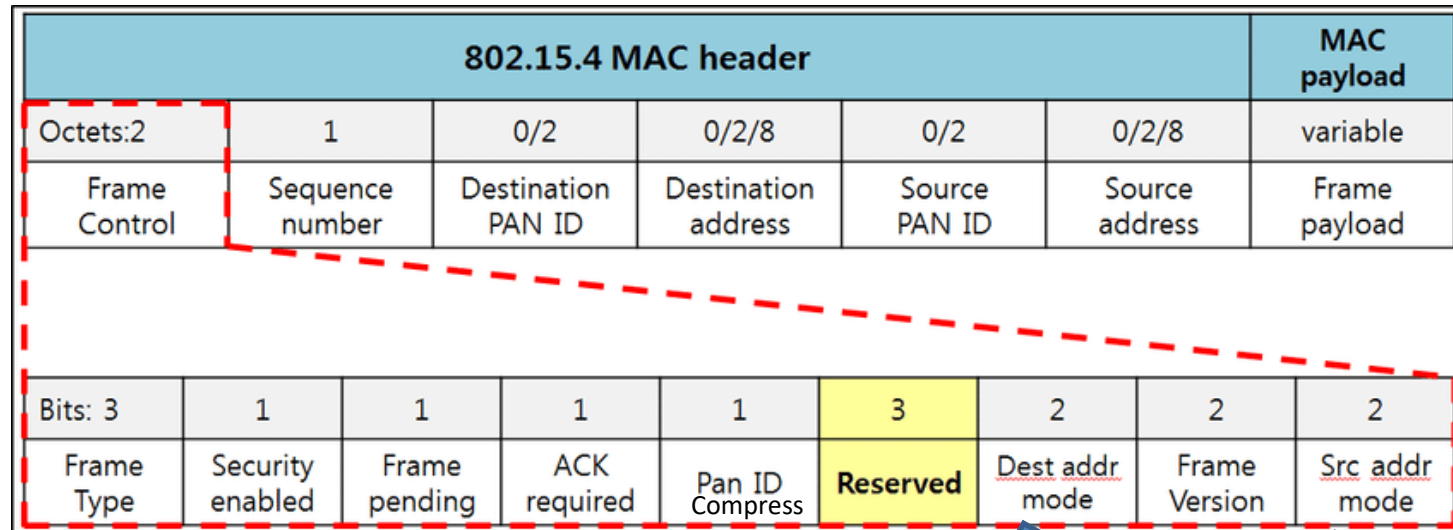
General MAC Frame Format



MAC frame types:

- Data frame
- ACK frame
- Beacon frame
- Command frame

Cont...

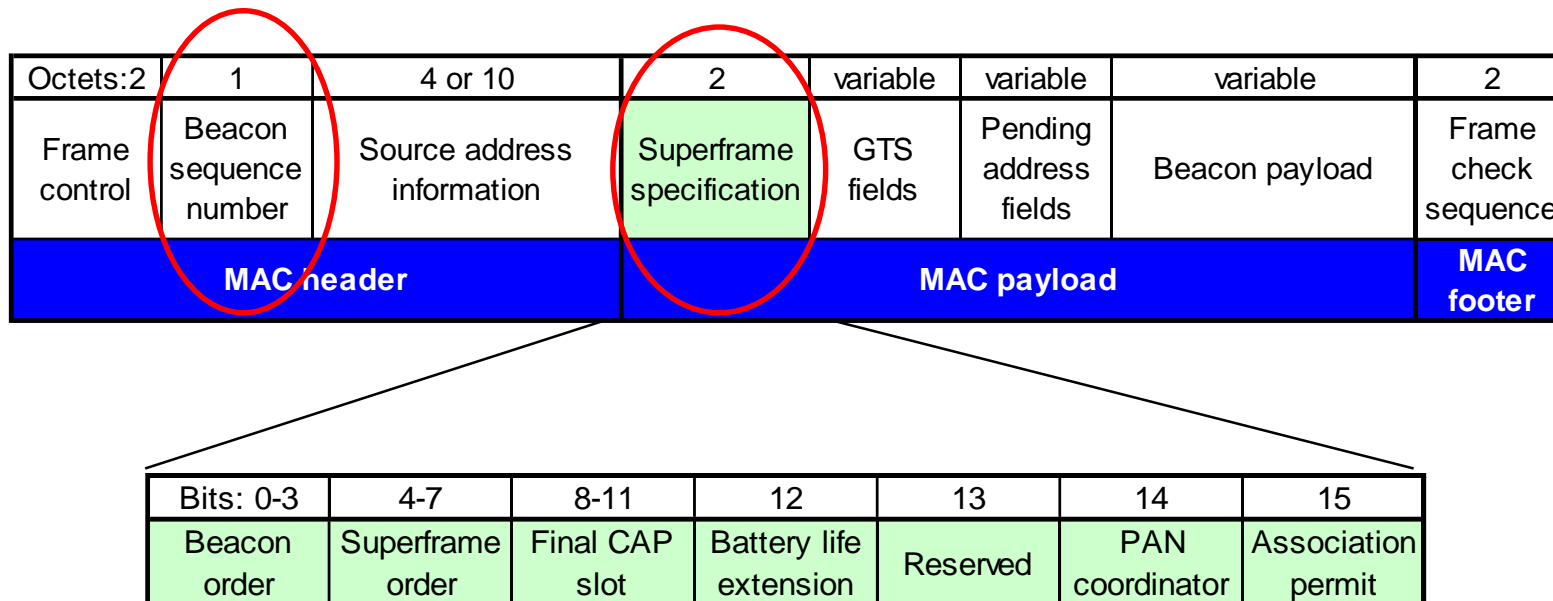


-Values of the Frame Type subfield

Frame type value $b_2 b_1 b_0$	Description
000	Beacon
001	Data
010	Acknowledgment
011	MAC command
100–111	Reserved

Addressing mode value $b_1 b_0$	Description
00	PAN identifier and address field are not present.
01	Reserved.
10	Address field contains a 16 bit short address.
11	Address field contains a 64 bit extended address.

Beacon Frame Format



Command Frame Format

Octets:2	1	4 to 20	1	variable	2
Frame control	Data sequence number	Address information	Command type	Command payload	Frame check sequence
MAC header			MAC payload		MAC footer

- Command Frame Types

- Association request
- Association response
- Disassociation notification
- Data request
- PAN ID conflict notification
- Orphan Notification
- Beacon request
- Coordinator realignment
- GTS request

Data & ACK Frame Format

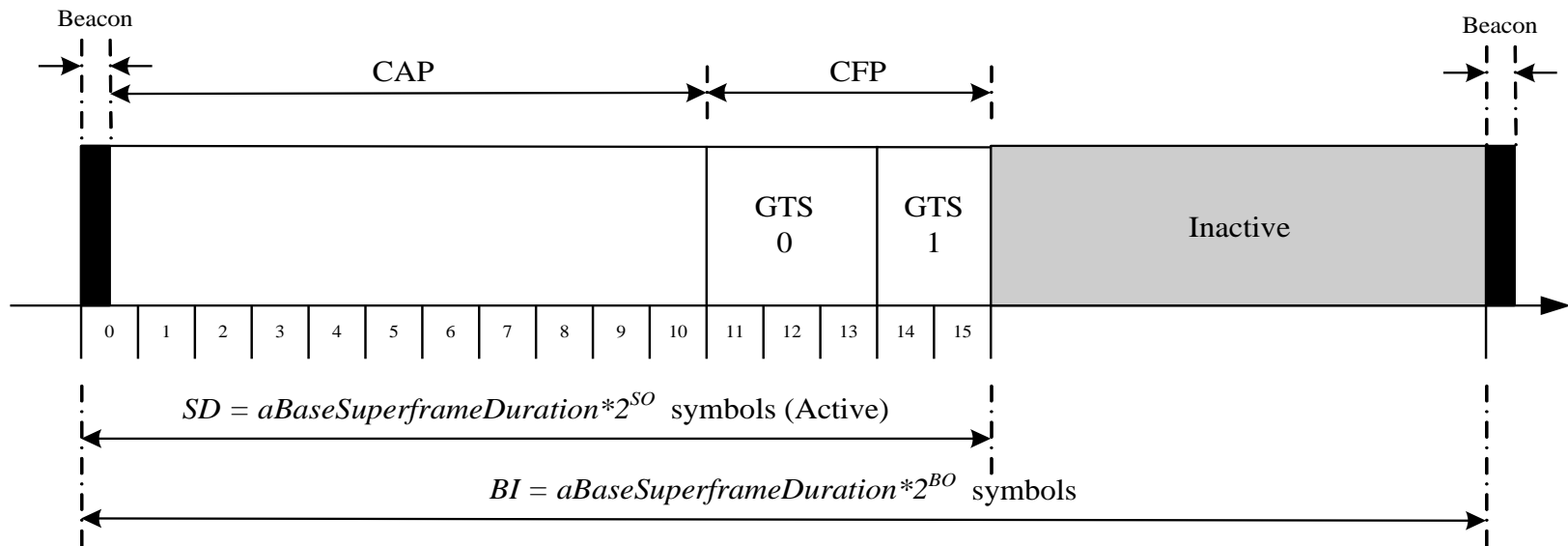
Data Frame

Octets:2	1	4 to 20	variable	2
Frame control	Data sequence number	Address information	Data payload	Frame check sequence
MAC header			MAC Payload	MAC footer

ACK Frame

Octets:2	1	2
Frame control	Data sequence number	Frame check sequence
MAC header		MAC footer

Superframe



- A superframe is divided into two parts

- **Inactive:** all station sleep.
 - no communication
 - nodes can turn their radios off and go into power saving mode
- **Active:**
 - Active period is divided into 16 slots in general
 - 16 slots are further divided into two parts
 - Contention access period (CAP)
 - Contention free period (CFP)
 - Beacon only period (BOP)

- **superframe order (SO)** : decides the length of the active portion in a superframe
- **beacon order (BO)** : decides the length of a superframe or beacon transmission period
- **beacon-enabled** network should satisfy $0 \leq SO \leq BO \leq 14$
- **PAN coordinator decides SO, BO**
 - Default value: SO=3, BO=5
- SD: Superframe Duration
- BI: Beacon Interval

Cont...



- ***aBaseSlotDuration***
 - = The number of symbols forming a superframe slot when *the superframe order (SO)* is equal to zero
 - = 60 PHY symbols
- ***aNumSuperframeSlots***
 - = The number of slots contained in any superframe
 - = 16
- ***aBaseSuperframeDuration***
 - = The number of symbols forming a superframe when *the superframe order (SO)* is equal to zero
 - = ***aBaseSlotDuration*** × ***aNumSuperframeSlots***
- So, Length of a superframe
 - = can range from 15.36 msec to 215.7 sec (= 3.5 min).
- Each device will be
 - active for $2^{-(BO-SO)}$ portion of the time
 - sleep for $1 - 2^{-(BO-SO)}$ portion of the time

Duty Cycle:	BO-SO	0	1	2	3	4	5	6	7	8	9	≥ 10
	Duty cycle (%)	100	50	25	12	6.25	3.125	1.56	0.78	0.39	0.195	< 0.1

Beacon Superframe and GTS

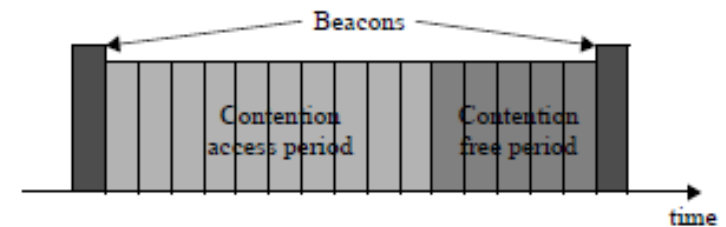
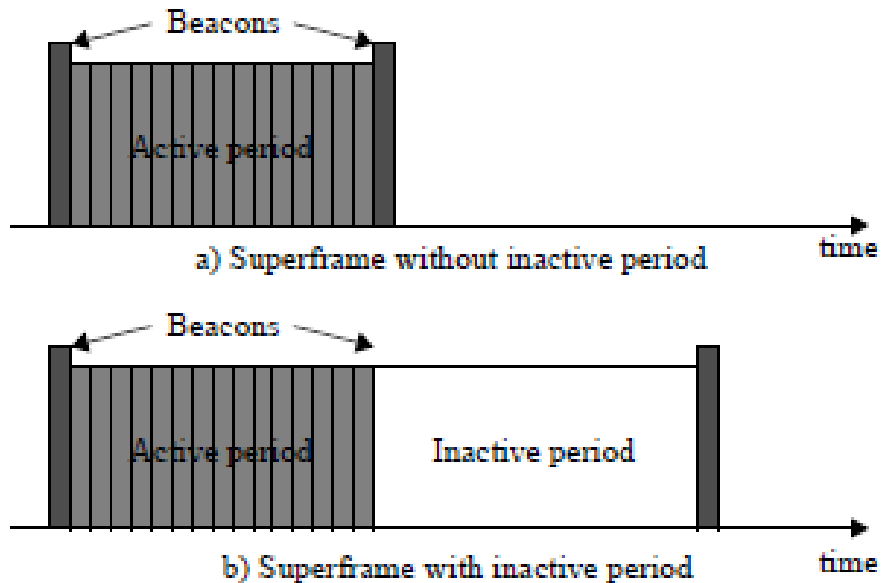


Figure 5-6—Structure of the active periods with GTSs

Active Periods with GTS

Beacon Enabled Superframe

Beaconed / Non-beaconed network

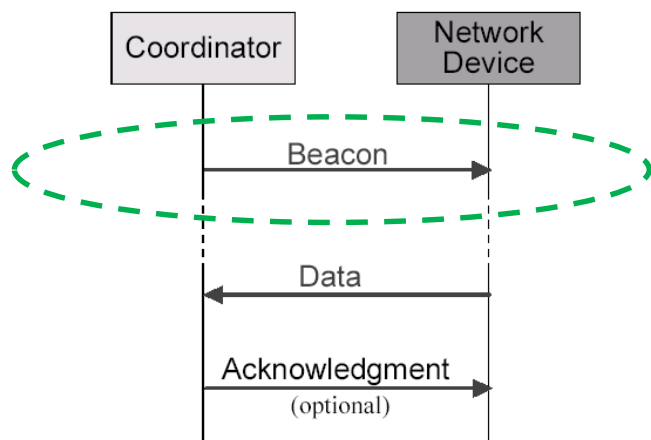


- In a “beacon-enabled” network (**i.e. uses superframe structure**)
 - Devices use the **slotted CAMA/CA** mechanism to contend for the channels
 - FFDs who require **fixed rates of transmissions** can ask for **GTS** from the coordinator
- In a “nonbeacon-enabled” network (**i.e. do not use superframe structure**)
 - Devices use the **unslotted CAMA/CA** mechanism for channel access
 - **GTS** shall not be permitted
- CSMA/CA is **not used for Beacon** transmission;
- CSMA/CA is also **not used** for **Data** transmission **during CFP**
- **Beacons** are used for
 - ❖ announcing the existence of a PAN
 - ❖ synchronizing with other devices
 - ❖ informing pending data in coordinators
 - ❖ starting superframes

Data Transfer: Device -> Coordinator

In a beacon-enabled network

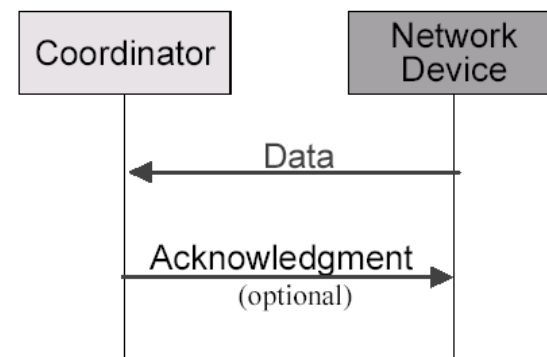
- a device finds the beacon to **synchronize** to the **superframe** structure.
- Then it uses **slotted CSMA/CA** to transmit its data.



Communication to a coordinator
In a **beacon-enabled** network

In a non-beacon-enabled network

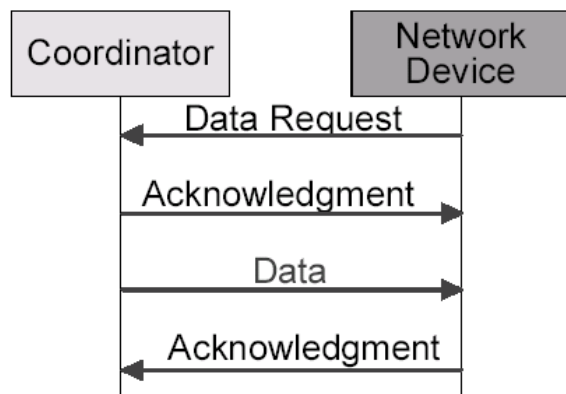
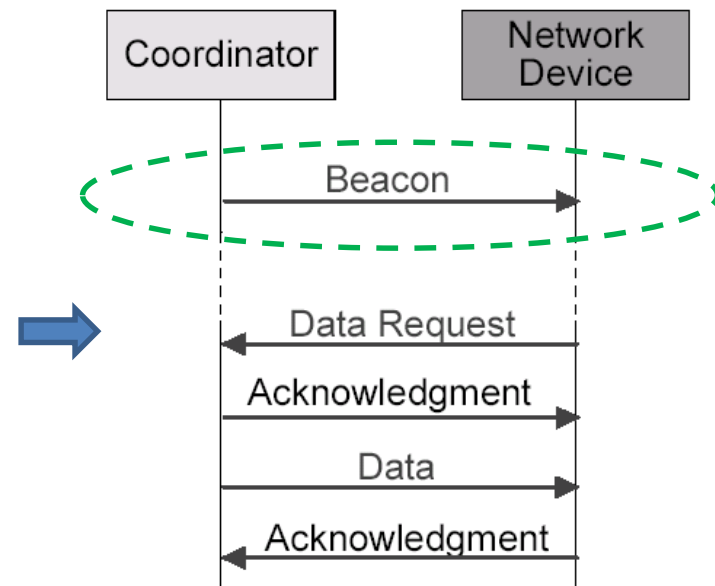
- device simply transmits its data using **unslotted CSMA/CA**



Communication to a coordinator
In a **non-beacon-enabled** network

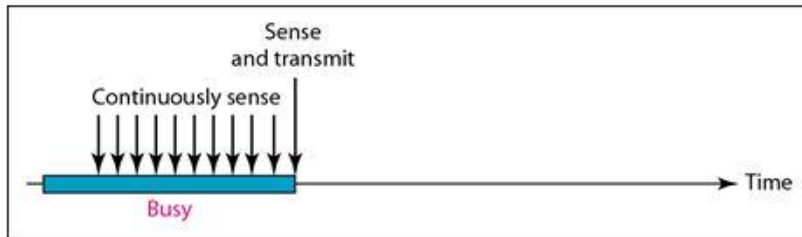
Data Transfer: Coordinator -> Device

- Data transferred **from coordinator to device**
 - in a **beacon-enabled** network:
 - The **coordinator indicates** in the **beacon** that some data is pending.
 - A device periodically listens to the beacon and transmits a **Data Request** command using **slotted CSMA/CA**.
 - Then **ACK, Data, and ACK**

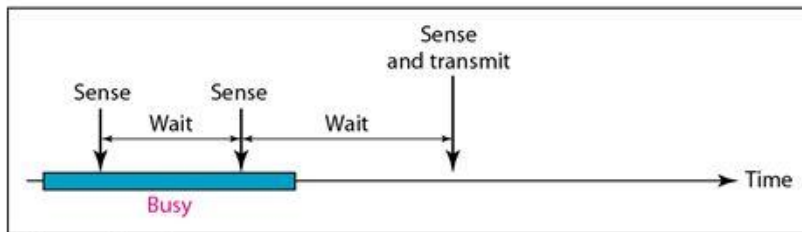


- Data transferred **from coordinator to device**
 - in a **non-beacon-enabled** network:
 - The device transmits a **Data Request** using **unslotted CSMA/CA**.
 - If the coordinator has its pending data, an **ACK** is replied.
 - Then the coordinator transmits **Data** using **unslotted CSMA/CA**.
 - If there is no pending data, a data frame with zero length payload is transmitted.
 - **ACK** is replied

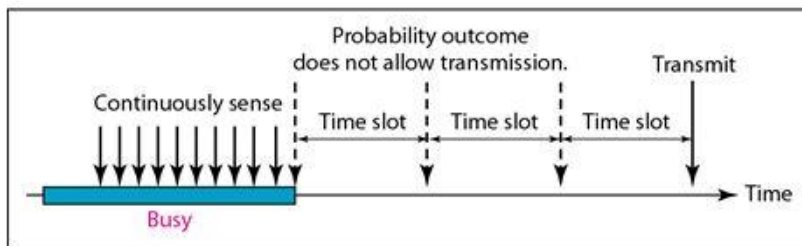
Channel Access Mechanism



a. 1-persistent



b. Nonpersistent



c. p-persistent

- **CSMA** requires that each station **first check the state of the medium** before sending.
- This method aims to **decrease the chances of collisions** when two or more stations want to transmit data
- **Persistent** methods can be applied to take action when the channel is sensed busy/idle.
 - **1-persistent**
 - When station found idle channel, it **transmits the frame without any delay**.
 - **Non-persistent**
 - when the channel is found busy, it will **wait for the random time** and again sense for the state of the station whether idle or busy
 - **p-persistent**
 - If the channel found to be idle, it **transmits the frame with probability p**
 - This is implemented using **backoff period** concept

Unslotted CSMA/CA

- CSMA/CA random channel access

- nonbeacon-enabled network → uses **unslotted CSMA/CA**

In unslotted CSMA/CA:

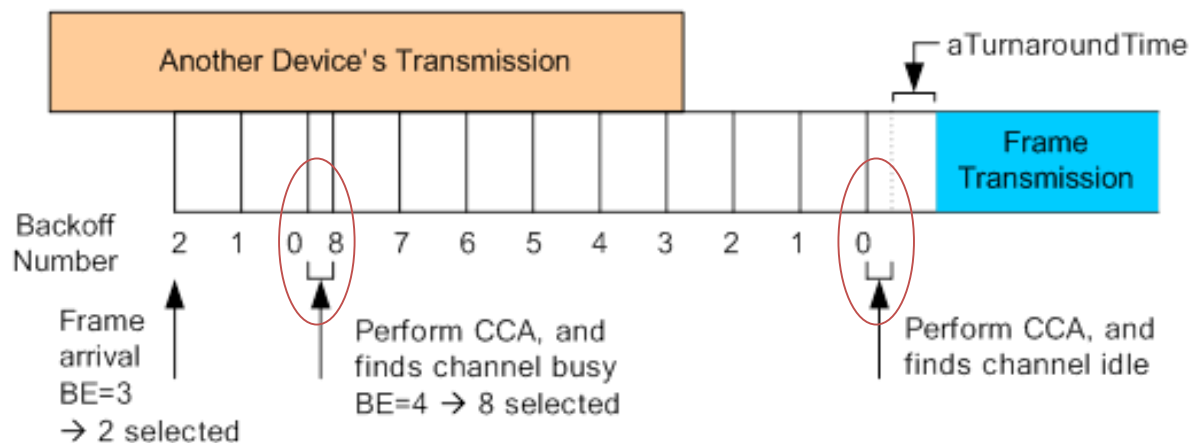
- The **backoff periods** of one device **are not related in time** to the backoff periods of any other device in the PAN.
- One backoff period = *aUnitBackoffPeriod*.

Backoff:

- is an algorithm that uses feedback to **multiplicatively decrease the rate** of some process

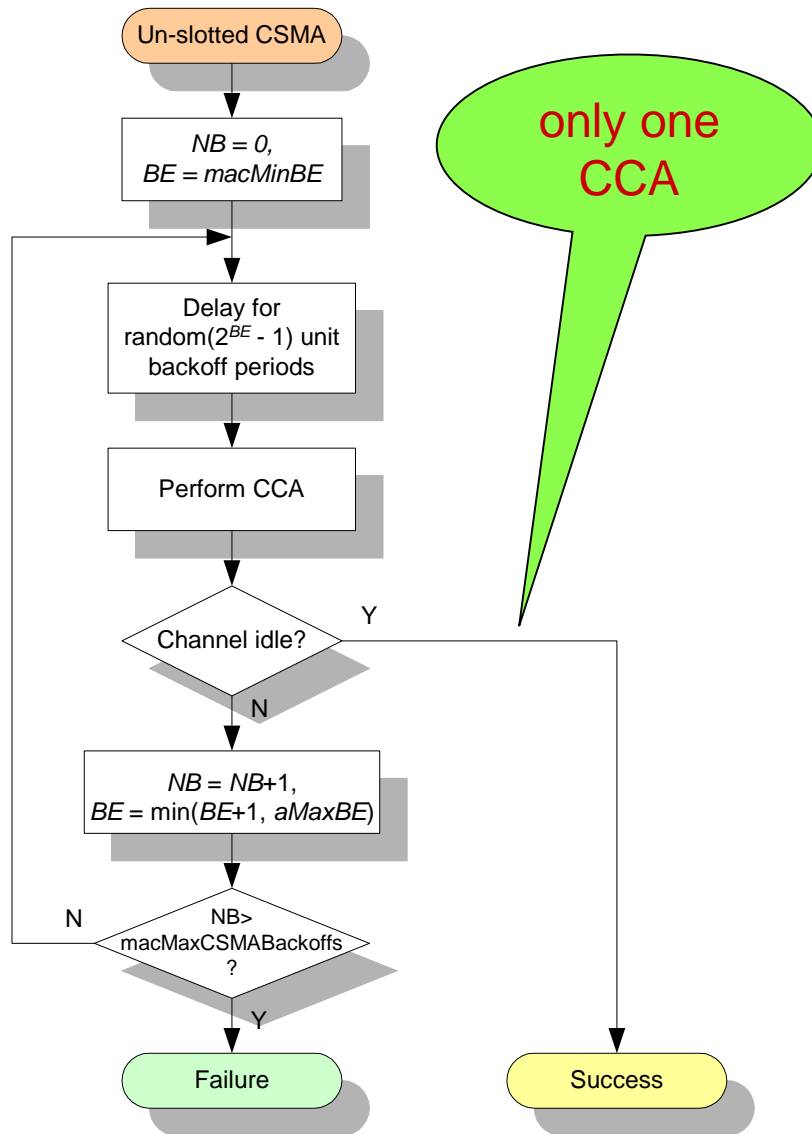
Binary exponential backoff (BEB)

- After **c collisions** in BEB algo., the delay is **randomly chosen** from $[0, 1, \dots, N]$ slots, where $N = 2^c - 1$.



BE: Backoff Exponent

Cont...



NB (Number of Backoff): number of times that backoff has been taken in this attempt of transmission

- if exceeding macMaxCSMABackoff , the attempt fails

BE (Backoff Exponent): play the role to decide how many backoff periods a device shall wait before attempting to assess a channel.

CCA (Clear Channel Assessment)

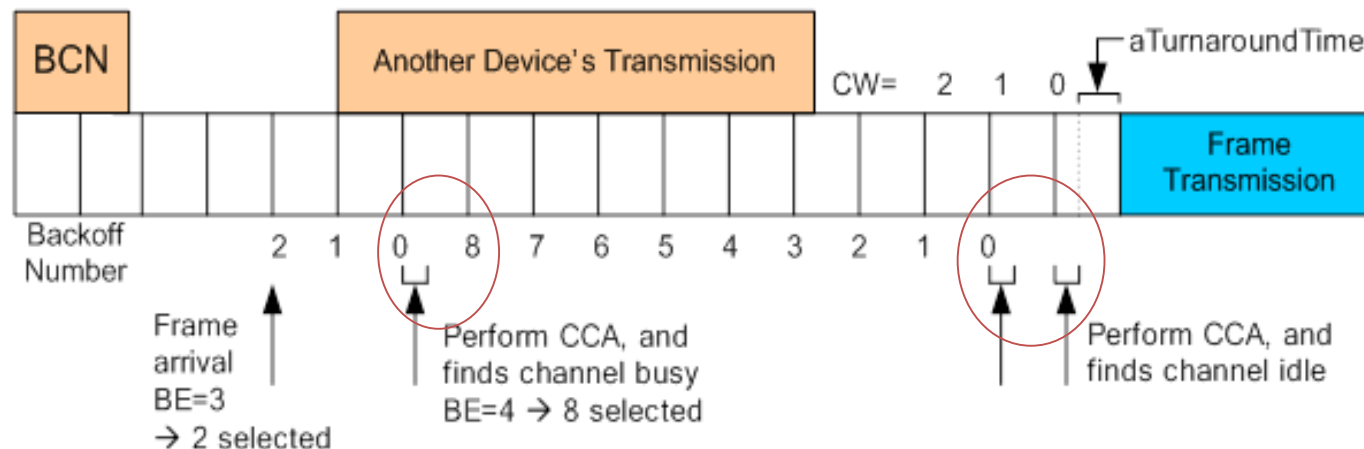
Slotted CSMA/CA

- CSMA/CA random channel access

➤ beacon-enabled network → uses **slotted CSMA/CA**

In slotted CSMA/CA:

- The **backoff period boundaries** of every **device** in the PAN shall be **aligned with** the superframe slot boundaries of the PAN coordinator
 - i.e. the **start of first backoff period** of each device is aligned with the **start of the beacon** transmission
- The MAC sublayer shall ensure that the PHY layer commences all of its **transmissions on the boundary of a backoff period**

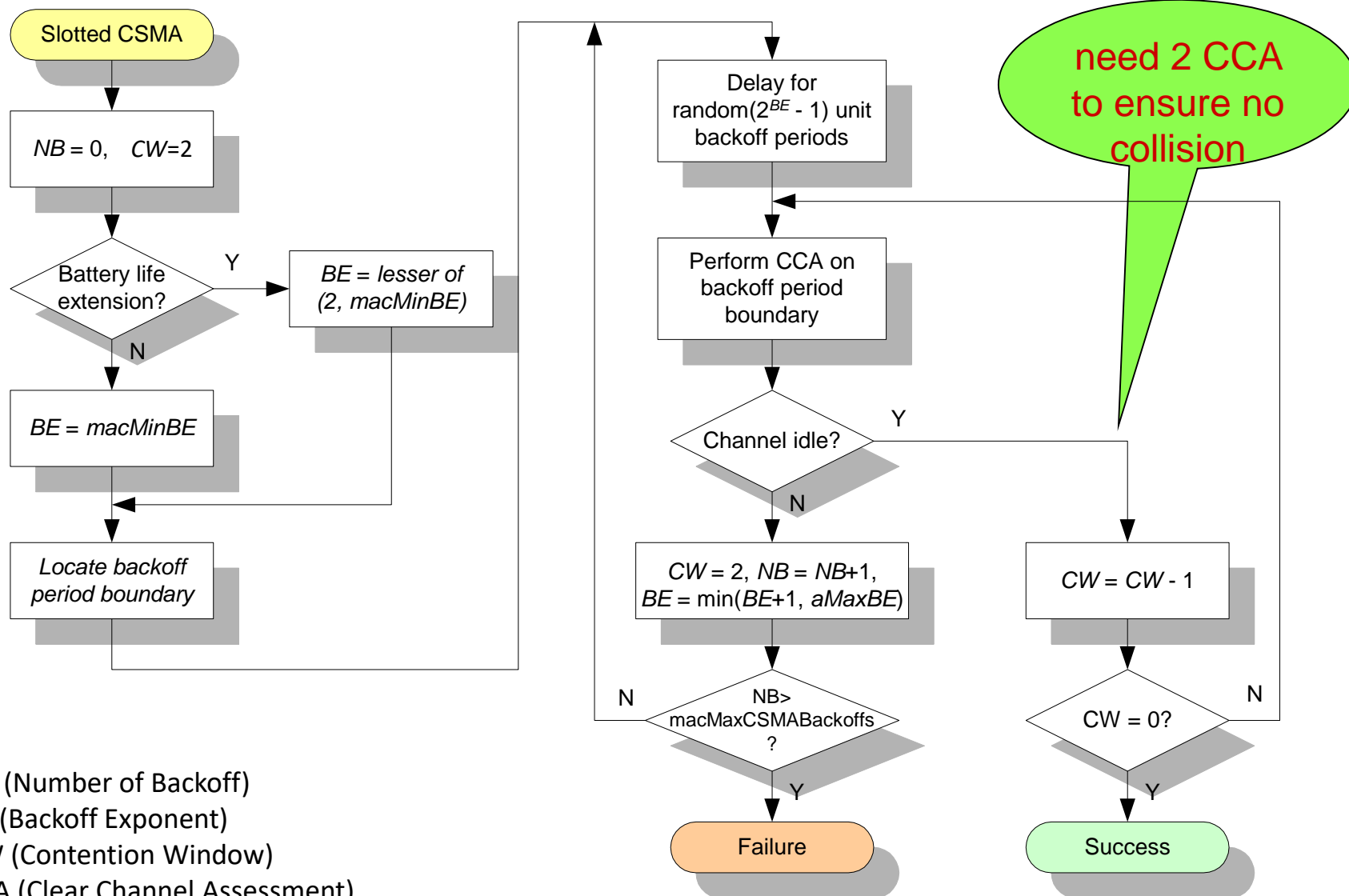


Cont...



- Each device maintains 3 variables for each transmission attempt
 - **NB (Number of Backoff)**: number of times that backoff has been taken in this attempt of transmission
 - if exceeding **macMaxCSMABackoff**, the attempt fails
 - **BE (Backoff Exponent)**: play the role to decide **how many backoff periods** a device shall wait before attempting to assess a channel.
 - The number of **backoff periods** is **lesser than** the remaining number of slots in the CAP
 - Otherwise, MAC sublayer shall **pause the backoff countdown at the end of the CAP**, and resume it at the start of the CAP in the next superframe
 - **CW (Contention Window)**: the number of clear slots that must be seen after each backoff
 - **always set to 2** and **count down to 0** if the channel is sensed to be clear
 - The design is for some PHY parameters, which require 2 CCA for efficient channel usage.
 - **Note**: CW in 802.15.4 is not same with CW in 802.11
 - CW in 802.11 is used to decide the backoff window size from which the backoff period is chosen randomly
 - CW in 802.15.4 is used to decide how many rounds of CCA is required before getting the channel access
- **Battery Life Extension (BLE)**:
 - designed for very low-power operation, where **a node only contends in the first few slots**

Cont...

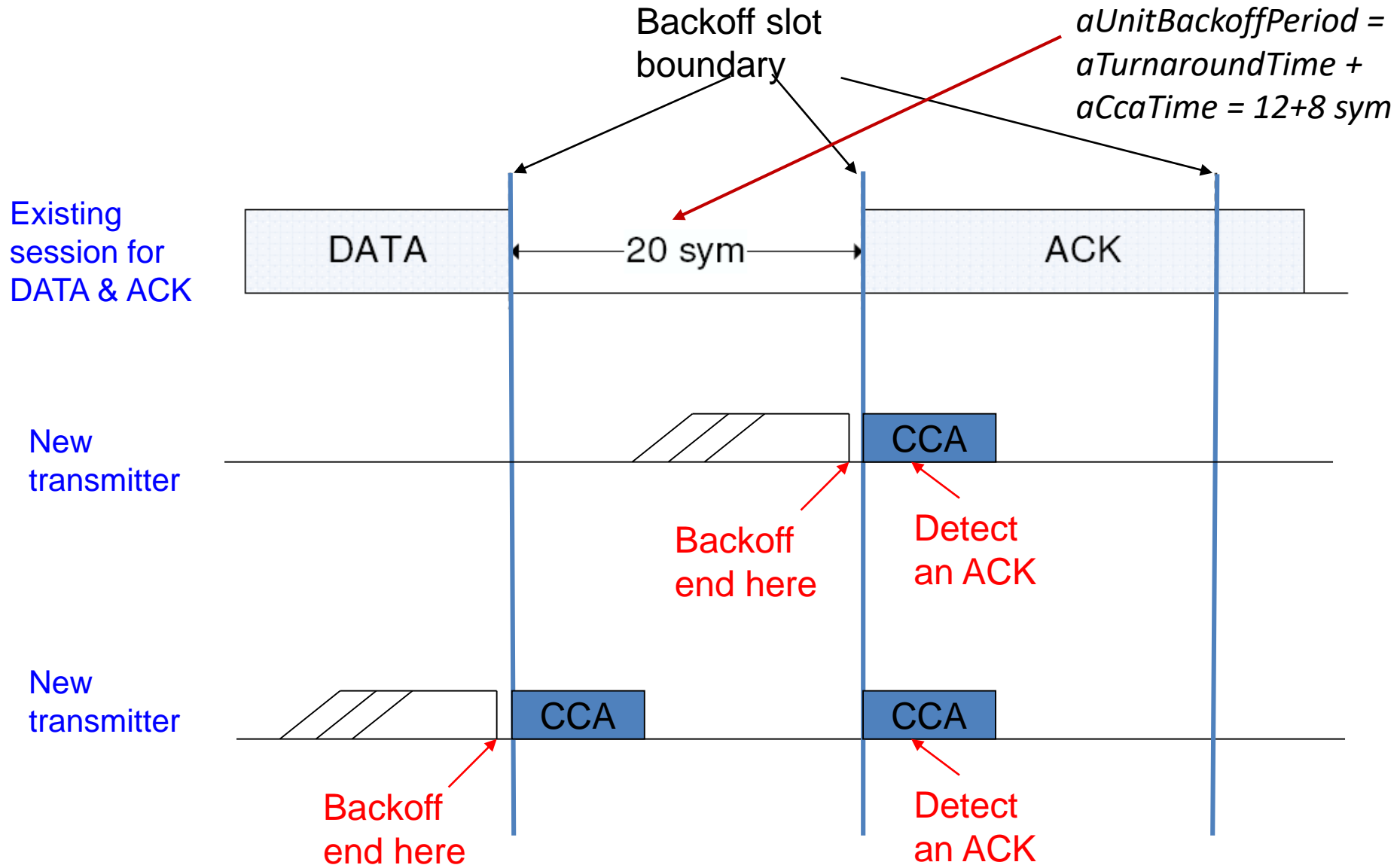


Why 2 CCAs to Ensure Collision-Free

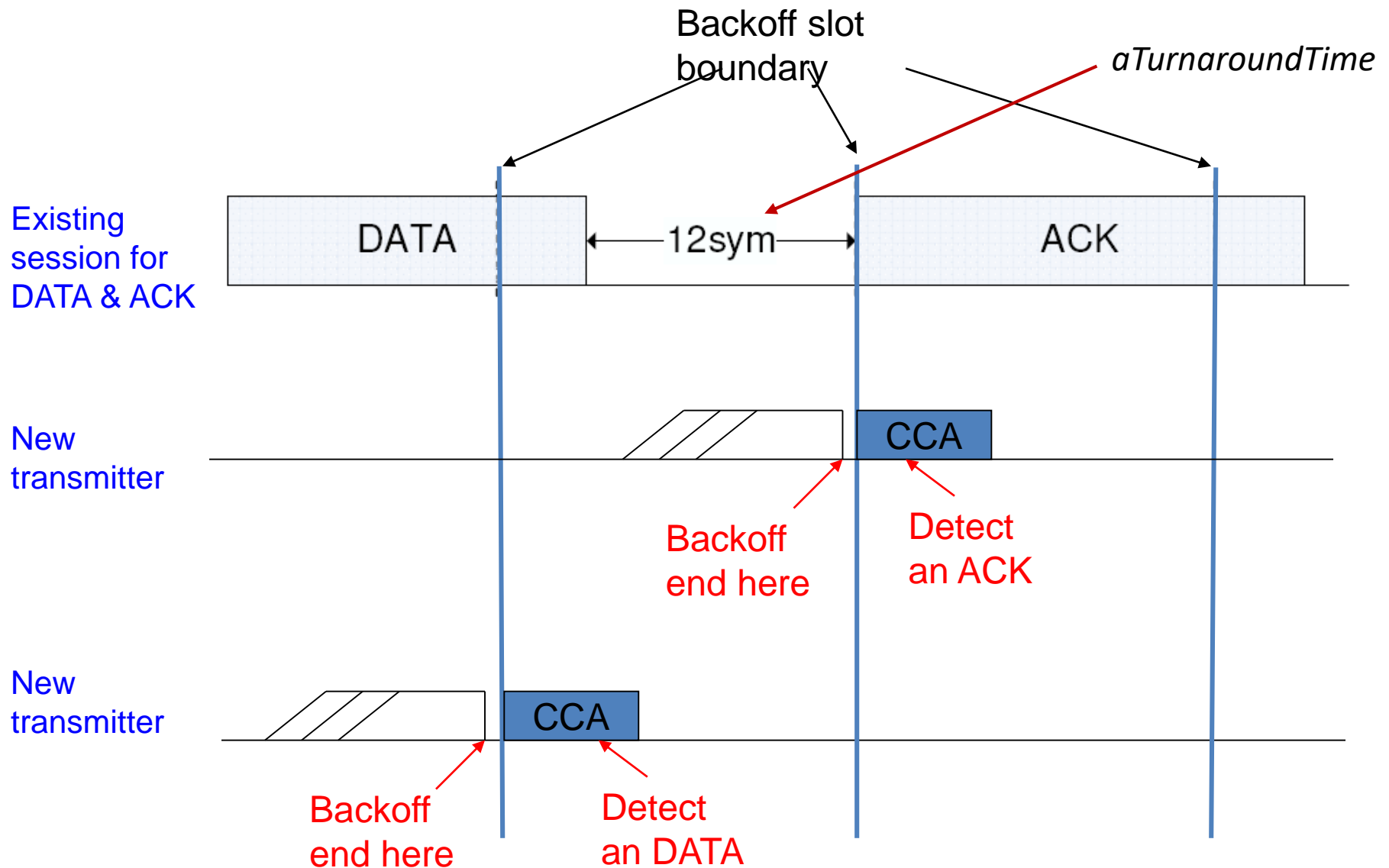


- Each CCA occurs at the boundary of a **backoff slot**
- Each **Backoff Slot duration** = 20 PHY symbols
- Each **CCA duration** = 8 PHY symbols
- The standard specifies that **a transmitter node performs the CCA twice in order to protect acknowledgment (ACK).**
 - When an ACK packet is expected, the receiver shall send it after a t_{ACK} time on the backoff boundary
 - t_{ACK} varies from 12 to 31 symbols
 - One-time CCA of a transmitter **may potentially cause a collision** between a **newly-transmitted packet** and an **ACK** packet.
 - (See examples below)

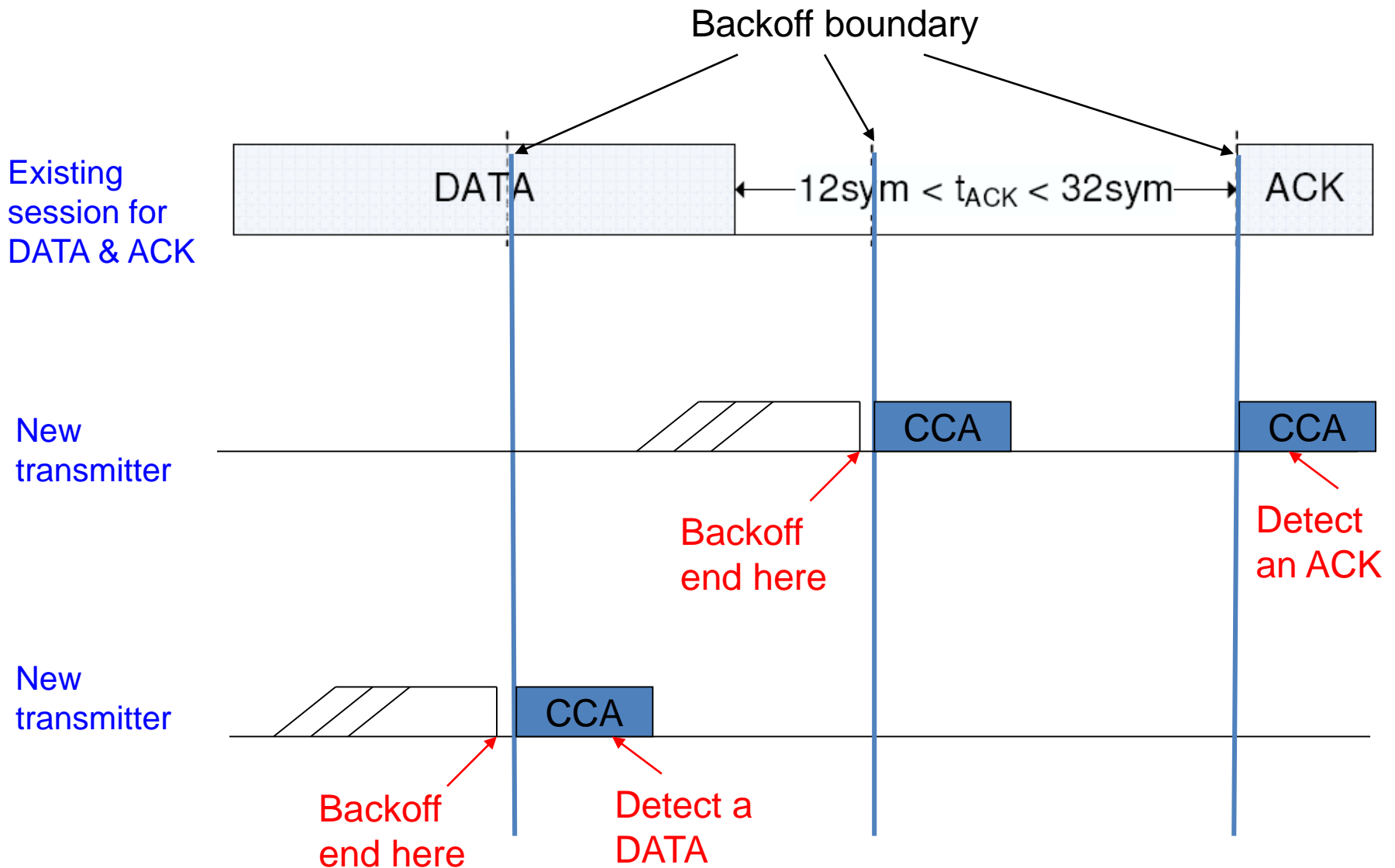
Why 2 CCAs (case 1)



Why 2 CCAs (Case 2)

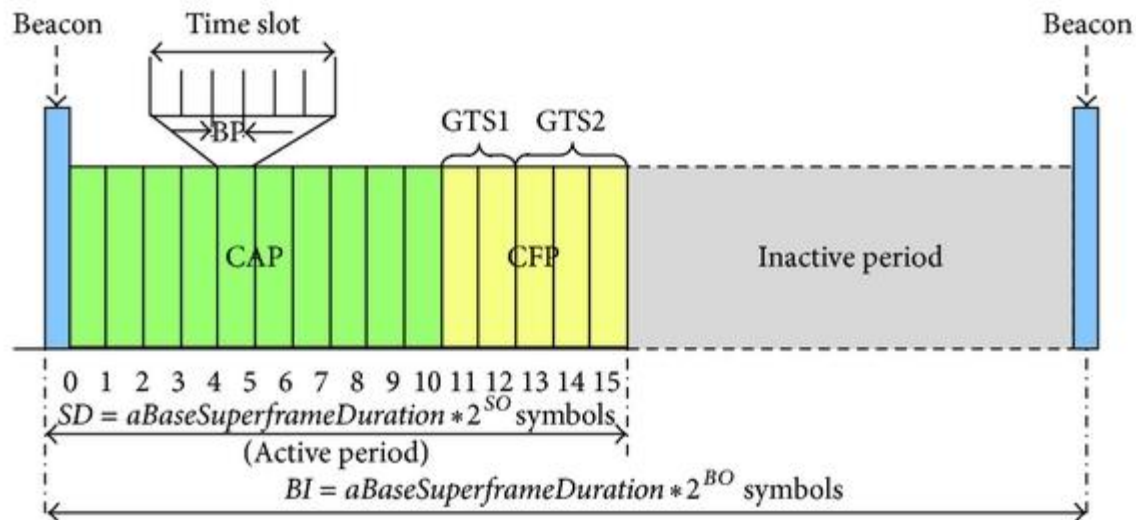


Why 2 CCAs (Case 3)



GTS Concepts

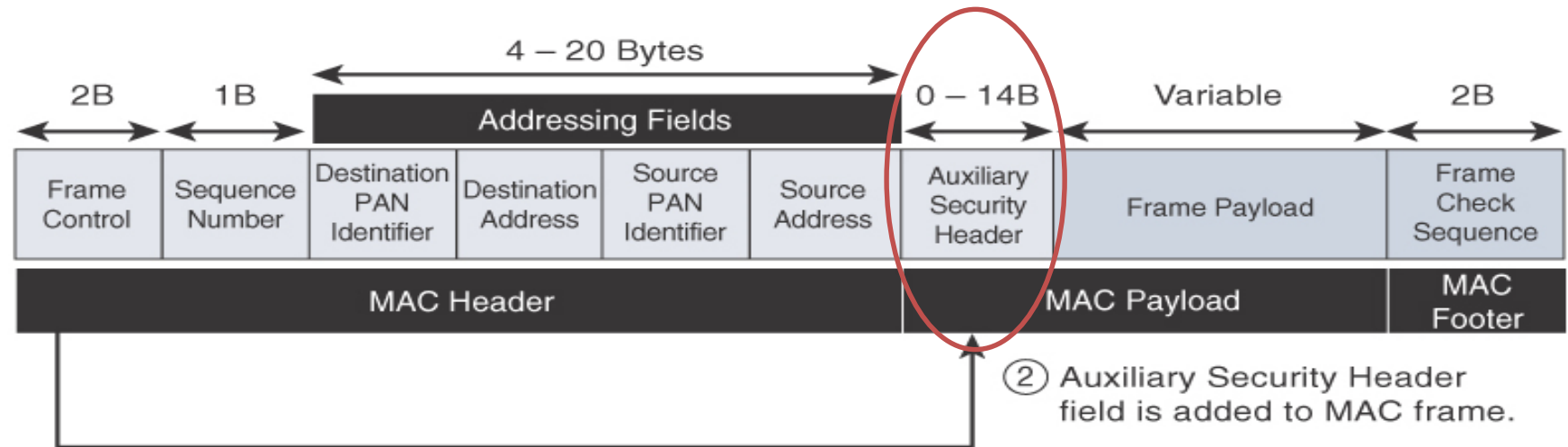
- A **guaranteed time slot (GTS)** allows a device to operate on the channel within a portion of the superframe
- A GTS shall only be allocated by the PAN coordinator
- The PAN coordinator can allocate up to **7 GTSs** at the same time
- The PAN coordinator decides whether to allocate GTS based on:
 - Requirements of the GTS request
 - The current available capacity in the superframe



Cont...

- A GTS can be deallocated
 - At any time at the discretion of the **PAN coordinator**, OR
 - **By the device** that originally requested the GTS
- A device that has been allocated a GTS may also operate in the CAP
- A data frame transmitted in an allocated GTS **shall use only short addressing**
- Before GTS starts, the **GTS direction** shall be specified as either Tx or Rx
 - Each device may request **one transmit GTS** and/or **one receive GTS**
- A device shall only attempt to allocate and use a GTS if it is currently tracking the beacon
- If a device **loses synchronization** with the PAN coordinator, all its **GTS allocations shall be lost**
- The use of GTSs by an RFD is optional

Security



① Security Enabled bit in Frame Control is set to 1.

- IEEE 802.15.4 specification uses **Advanced Encryption Standard (AES)** with a 128-bit key length as the base encryption algorithm
- Message integrity code (MIC)**, which is calculated for the entire frame using the same AES key, to validate the data that is sent

Limitations in 802.15.4



- **Disadvantages of Initial version of IEEE 802.15.4**
 - MAC reliability
 - unbounded latency
 - multipath fading

- **IEEE 802.15.4e** amendment of IEEE 802.15.4-2011 expands the MAC layer feature set

- to remedy the disadvantages of 802.15.4.
- to better suitable in factory and process automation, and smart grid
- **Main modifications** were:
 - frame format,
 - security,
 - determinism mechanism,
 - frequency hopping

- **IEEE 802.15.4g** amendment of IEEE 802.15.4-2011 expands the PHY layer feature set

- to optimize large outdoor wireless mesh networks for field area networks (FANs)
- to better suitable in smart grid or smart utility network (SUN) communication
- **Main modifications** were:
 - New PHY definitions
 - some MAC modifications were needed to support the new PHY

Lessons Learned



- ✓ IEEE 802.15.4 MAC
 - Network Topology
 - Device Types
 - Device Addressing
 - MAC Frame Formats

 - Timeslot, Superframe
 - Data Transfer Model

 - Channel Access Methods
 - Slotted CSMA/CA

 - Guaranteed time slot (GTS)
 - Association Procedure
 - Security

- ✓ Limitations of IEEE 802.15.4

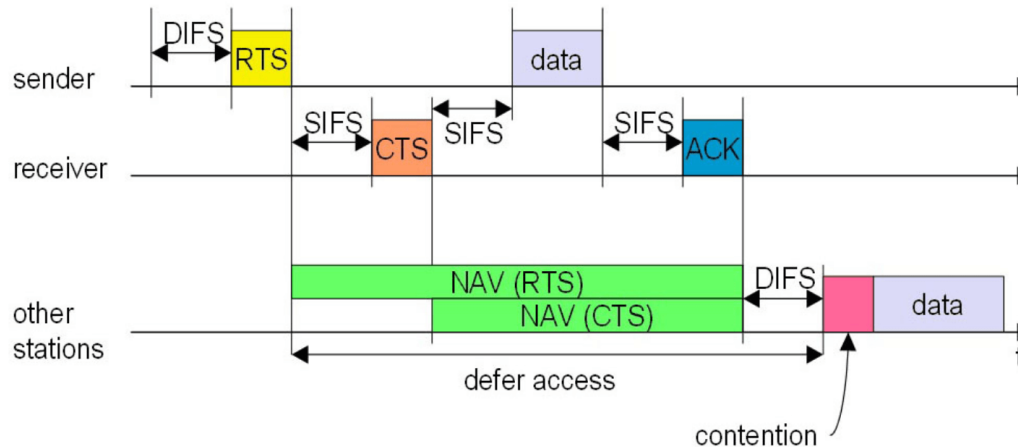
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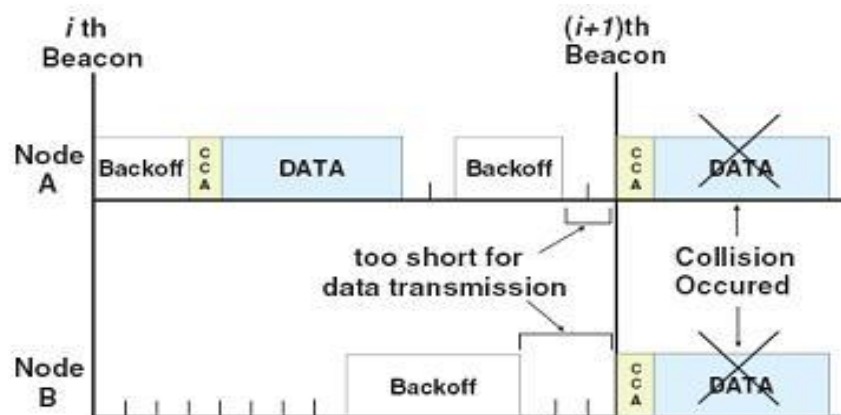
Figures and slide materials are taken from the following sources:

1. David Hanes *et al.*, “IoT Fundamentals: Networking Technologies, Protocols, and Use Cases for the Internet of Things”, 1st Edition, 2018, Pearson India.
2. Oliver Hersent et al., “The Internet of Things: Key Applications and Protocols”, 2018, Wiley India Pvt. Ltd.

Contention in 802.11 & 802.15.4



Contention in
IEEE 802.11 DCF



Contention in
IEEE 802.15.4
(for slotted CSMA/CA)